

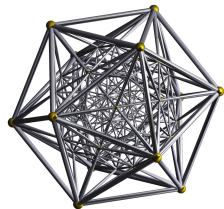
# Computational Principles for High-dim Data Analysis

## (Lecture Two)

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# Sparse Signal Models

## ① Applications

- Medical Imaging
- Image Processing
- Face Recognition

## ② Recovering a Sparse Signal

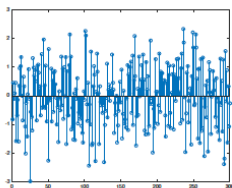
- Norms for Sparsity
- The  $\ell^0$  Norm
- Minimizing the  $\ell^0$  for the Sparsest Solution
- Computational Complexity of  $\ell^0$  Minimization

# Recovering Signals from (Linear) Measurements

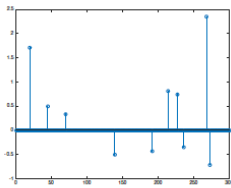
$$\underset{\text{observation}}{\mathbf{y}} = \underset{\text{measurement matrix}}{\mathbf{A}} \underset{\text{unknown}}{\mathbf{x}}. \quad (1)$$

$$\mathbf{y} \in \mathbb{R}^m, \quad \mathbf{x} \in \mathbb{R}^n, \quad \mathbf{A} \in \mathbb{R}^{m \times n}, m \ll n.$$

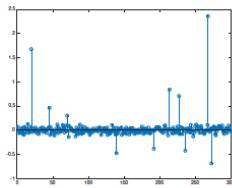
Different “structures” in  $\mathbf{x}$ :



dense vector



sparse vector



compressible vector

# Applications

# Application I: Magnetic Resonance Imaging (Chap. 10)

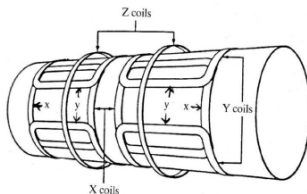
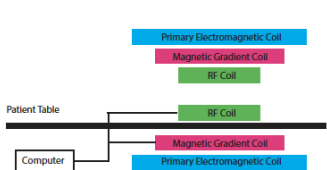
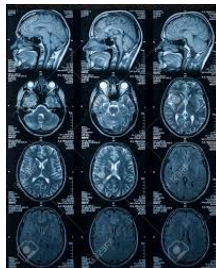


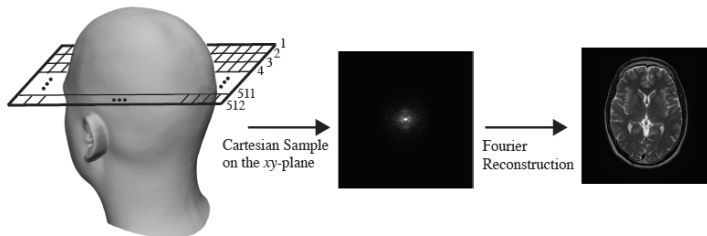
Figure: Left: Key components. Right: The three-axis gradient coils.

# Mathematical Model of MRI

Simplified mathematical model for MRI:

$$y = \mathcal{F}[I](\mathbf{u}) = \int_{\mathbf{v}} I(\mathbf{v}) \exp(-i 2\pi \mathbf{u}^* \mathbf{v}) d\mathbf{v}, \quad \mathbf{u}, \mathbf{v} \in \mathbb{R}^2 \quad (2)$$

$$\mathbf{y} = \begin{bmatrix} y_1 \\ \vdots \\ y_m \end{bmatrix} = \begin{bmatrix} \mathcal{F}[I](\mathbf{u}_1) \\ \vdots \\ \mathcal{F}[I](\mathbf{u}_m) \end{bmatrix} \doteq \mathcal{F}_U[I], \quad m \ll N^2. \quad (3)$$



**Figure:** Recovering MRI image from Fourier measurements.

# Exploit Structures of MRI Images

Express  $I$  as a superposition of basis functions  $\Psi = \{\psi_1, \dots, \psi_{N^2}\}$ :

$$\underset{\text{image}}{I} = \sum_{i=1}^{N^2} \underset{i\text{-th basis signal}}{\psi_i} \times \underset{i\text{-th coefficient}}{x_i}. \quad (4)$$

Approximate  $I$  with the  $k$  largest coefficients  $J = \{i_1, \dots, i_k\}$ :

$$\underset{\text{target image}}{I} \approx \underset{\text{superposition of } k \text{ basis functions}}{\tilde{I}_k = \sum_{i \in J} \psi_i x_i}. \quad (5)$$

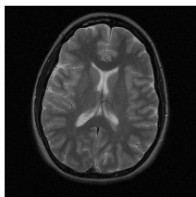
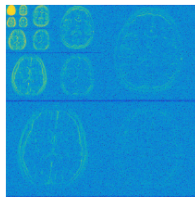


image  $I(v)$



wavelet coefficients  $x: I = \Psi[x]$ .

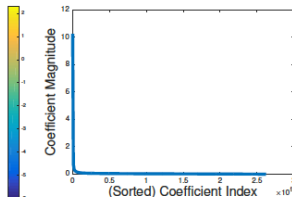


Figure: MRI image and its wavelet coefficients

# Recovering Image from a Under-determined Linear System

$$\begin{aligned}
 \text{observed Fourier coefficients } \mathbf{y} &= \mathcal{F}_U[I], \\
 &= \mathcal{F}_U[\psi_1 x_1 + \cdots + \psi_{N^2} x_{N^2}], \\
 &= \mathcal{F}_U[\psi_1] x_1 + \cdots + \mathcal{F}_U[\psi_{N^2}] x_{N^2}, \\
 &= [\mathcal{F}_U[\psi_1] \mid \cdots \mid \mathcal{F}_U[\psi_{N^2}]] \mathbf{x}, \\
 &\quad \text{matrix } \mathbf{A} \in \mathbb{R}^{m \times N^2}, m \ll N^2. \\
 &= \mathbf{A} \mathbf{x}.
 \end{aligned} \tag{6}$$

**$\mathbf{x}$  is sparse or approximately sparse!**

$$\mathbf{I} \approx \sum_{i \in J} \psi_i x_i = \underbrace{\Psi}_{N^2 \times N^2 \text{ matrix}} \underbrace{\mathbf{x}}_{\text{sparse vector}}, \tag{7}$$

where  $x_i = 0$  for  $i \notin J$ , and  $k = |J| \ll N^2$ .



## Application II: Image Denoising

$$I_{\text{noisy}} = \underbrace{I_{\text{clean}}}_{\text{target image}} + \underbrace{z}_{\text{noise}} \quad (8)$$

Break  $I_{\text{clean}}$  into patches  $y_{1\text{clean}}, \dots, y_{p\text{clean}}$ :

$$y_i = y_{i\text{clean}} + z_i = \underbrace{A}_{\text{patch dictionary}} \times \underbrace{x_i}_{\text{sparse coefficient vector}} + z_i \quad (9)$$



**Figure:** Left: input image; middle: denoised; right: dictionary for image patches.

*Sparse representation for color image restoration*, Mairal, Elad, and Sapiro. TIP, 2008

## Application II: Image Super-Resolution

Given a pair of corresponding low-resolution and high-resolution dictionaries ( $\mathbf{A}_{\text{low}}, \mathbf{A}_{\text{high}}$ ):

$$\mathbf{y}_{i\text{low}} = \mathbf{A}_{\text{low}}\mathbf{x}_i \xrightarrow{\text{Lifting}} \mathbf{y}_{i\text{high}} = \mathbf{A}_{\text{high}}\mathbf{x}_i. \quad (10)$$

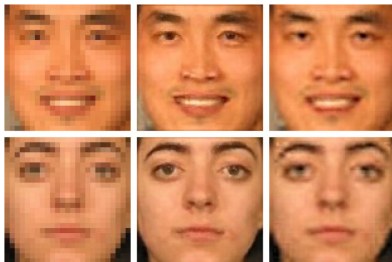
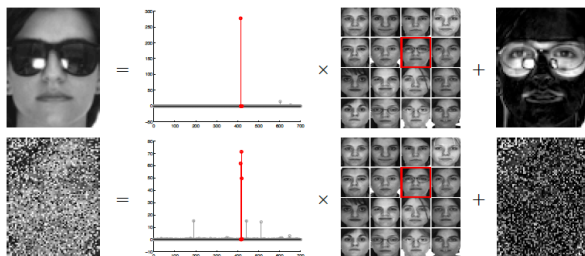


Figure: Super resolution for face images

*Image Super-Resolution via Sparse Representation of Raw Image Patches*, Yang, Wright, Huang, and Yi Ma, CVPR, 2008.

# Application III: Robust Face Recognition (Chap. 13)



$$\underset{\text{observation}}{\mathbf{y}} = \underset{\text{clean data}}{\mathbf{y}_o} + \underset{\text{sparse error}}{\mathbf{e}} \in \mathbb{R}^m. \quad (11)$$

Concatenate gallery images of  $n$  subjects into a large “dictionary”:

$$\underset{\text{all training images}}{\mathbf{B}} = [\mathbf{B}_1 \mid \mathbf{B}_2 \mid \cdots \mid \mathbf{B}_n] \in \mathbb{R}^{m \times n}, \quad n = \sum_i n_i. \quad (12)$$

# Application III: Robust Face Recognition (Chap. 13)

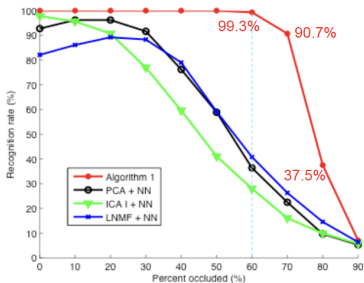
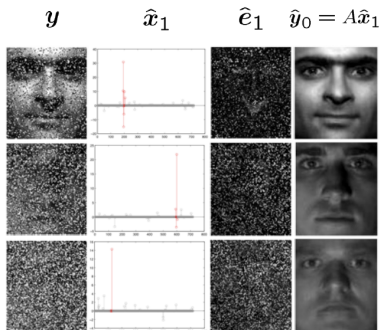
Find sparse solutions  $(x, e)$  to the linear system:

$$y = Bx + e = [B, I] \begin{bmatrix} x \\ e \end{bmatrix}. \quad (13)$$

Extended Yale B Database  
(38 subjects)

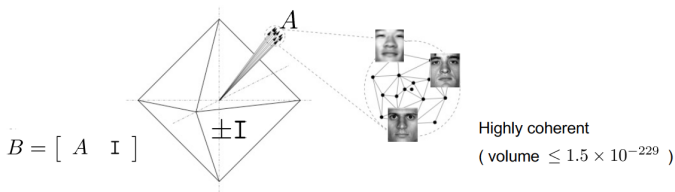
Training: subsets 1 and 2 (717 images)

Testing: subset 3 (453 images)



*Robust Face Recognition via Sparse Representation*, Wright, Yang, Ganesh, Sastry, and Ma, TPAMI, 2009.

# Robust Face Recognition: Rigorous Justification



**Theorem 1.** For any  $\delta > 0$ ,  $\exists \nu_0(\delta) > 0$  such that if  $\nu < \nu_0$  and  $\rho < 1$ , in weak proportional growth, with error support  $J$  and signs  $\sigma$  chosen uniformly at random,

$$\lim_{m \rightarrow \infty} P_{A,J,\sigma} \left[ \ell^1\text{-recoverability at } (I, J, \sigma) \ \forall I \in \binom{[n]}{k_1} \right] = 1.$$

*“ $\ell^1$  recovers any sparse signal from almost any error with density less than 1”*

*Dense Error Correction via  $\ell^1$  Minimization*, Wright and Ma, Trans. Info. Theory, 2010.

# Many Many Applications...

## Part III of the Textbook:

- Magnetic Resonance Imaging (Chapter 10)
- Wideband Spectrum Sensing (Chapter 11)
- Scientific Imaging Problems (Chapter 12)
- Robust Face Recognition (Chapter 13)
- Robust Photometric Stereo (Chapter 14)
- Structured Texture Recovery (Chapter 15)
- Deep Networks for Classification (Chapter 16)

**Your final projects?**

# How to Find a Sparse Solution $x$ ?

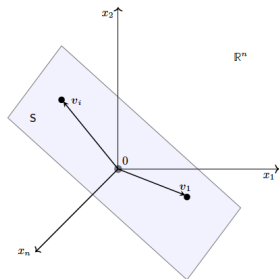
$$y = \underset{m \times n, m \ll n}{A} \times \underset{k \text{ sparse}}{x}. \quad (14)$$

$$y \in \mathbb{R}^m = A x \in \mathbb{R}^n$$

# Vector Spaces and Linear Algebra

## Topics in **Linear Algebra** (Appendix A):

- Vector Space, Linear Independence, Basis
- Linear Mappings, Subspaces, Matrix Representation
- Linear Systems and Conjugate Gradient (over-determined, under-determined)
- Eigenvalue Decomposition and Singular Value Decomposition
- Norms and Matrix Norms





# Norms

A *norm* on a vector space  $\mathbb{V}$  over  $\mathbb{R}$  is a function  $\|\cdot\| : \mathbb{V} \rightarrow \mathbb{R}$  that is

- ① **positive definite:**  $\|\mathbf{x}\| \geq 0$ , and  $\|\mathbf{x}\| = 0$  if and only if  $\mathbf{x} = \mathbf{0}$ ;
- ② **nonnegatively homogeneous:**

$$\|\alpha \mathbf{x}\| = |\alpha| \|\mathbf{x}\|, \quad \forall \mathbf{x} \in \mathbb{V}, \alpha \in \mathbb{R};$$

- ③ **subadditive:**  $\|\cdot\|$  satisfies the triangle inequality

$$\|\mathbf{x} + \mathbf{y}\| \leq \|\mathbf{x}\| + \|\mathbf{y}\|, \quad \forall \mathbf{x}, \mathbf{y} \in \mathbb{V}.$$

**Example:** for  $\mathbb{V} = (\mathbb{R}^n, \mathbb{R})$ , a  $p$ -norm:

$$\|\mathbf{x}\|_p \doteq \left( \sum_i |x_i|^p \right)^{1/p}, \quad p \in [1, \infty]. \quad (15)$$

# Sparse Vectors and Norms

$\ell^\infty$  norm or the “cube” norm:

$$\|\mathbf{x}\|_\infty = \max_i |x_i|. \quad (16)$$

$\ell^2$  norm or the “Euclidean norm”:

$$\|\mathbf{x}\|_2 = \sqrt{\sum_i |x_i|^2} = \sqrt{\mathbf{x}^* \mathbf{x}}. \quad (17)$$

$\ell^1$  norm:

$$\|\mathbf{x}\|_1 = \sum_i |x_i|, \quad (18)$$

$\ell^p$  unit ball:

$$\mathbf{B}_p \doteq \{\mathbf{x} \mid \|\mathbf{x}\|_p \leq 1\} \quad 0 < p \leq \infty. \quad (19)$$

**Example:** relative volumes of the  $\ell^p$  balls when  $n$  is large.

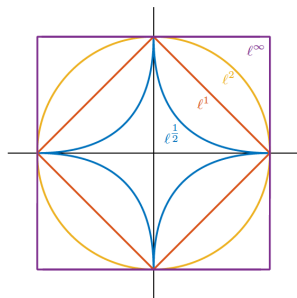


Figure: The  $\ell^p$  balls.

# Measure of Sparsity: the $\ell^0$ Norm

$\ell^0$  “norm”:

$$\|\mathbf{x}\|_0 = \#\{i \mid \mathbf{x}(i) \neq 0\}. \quad (20)$$

Not a norm, and only subadditive:

$$\forall \mathbf{x}, \mathbf{x}', \quad \|\mathbf{x} + \mathbf{x}'\|_0 \leq \|\mathbf{x}\|_0 + \|\mathbf{x}'\|_0. \quad (21)$$

“Limit” of  $\ell^p$ -norm as  $p \searrow 0$ :

$$\lim_{p \searrow 0} \|\mathbf{x}\|_p^p = \sum_{i=1}^n \lim_{p \searrow 0} |\mathbf{x}(i)|^p = \sum_{i=1}^n \mathbb{1}_{\mathbf{x}(i) \neq 0} = \|\mathbf{x}\|_0. \quad (22)$$

# Minimizing the $\ell^0$ Norm

Given  $\mathbf{y} = \mathbf{A}\mathbf{x}_o$ , to find  $\mathbf{x}_o$  as the sparsest solution to  $\mathbf{y} = \mathbf{A}\mathbf{x}$ , we can:

$$\hat{\mathbf{x}} = \arg \min \|\mathbf{x}\|_0 \quad \text{s.t.} \quad \mathbf{A}\mathbf{x} = \mathbf{y}. \quad (23)$$

The “support” of the nonzero entries:

$$\text{supp}(\mathbf{x}) = \{i \mid \mathbf{x}(i) \neq 0\} \subset \{1, \dots, n\}, \quad k = |\text{supp}(\mathbf{x}_o)|. \quad (24)$$

Brute force exhaustive search:

$$\mathbf{A}_I \mathbf{x}_I = \mathbf{y} \quad \forall I \subseteq \{1, \dots, n\}, \quad |I| \leq k. \quad (25)$$

# Uniqueness of the Sparsest Solution

## Definition (Kruskal Rank)

The *Kruskal rank* of a matrix  $\mathbf{A}$ , written as  $\text{krank}(\mathbf{A})$ , is the largest number  $r$  such that every subset of  $r$  columns of  $\mathbf{A}$  is linearly independent.

## Theorem ( $\ell^0$ Recovery)

Suppose that  $\mathbf{y} = \mathbf{A}\mathbf{x}_o$ , with

$$\|\mathbf{x}_o\|_0 \leq \frac{1}{2} \text{krank}(\mathbf{A}). \quad (26)$$

Then  $\mathbf{x}_o$  is the unique optimal solution to the  $\ell^0$  minimization problem

$$\min \|\mathbf{x}\|_0 \quad \text{s.t.} \quad \mathbf{A}\mathbf{x} = \mathbf{y}. \quad (27)$$

**Proof:**  $\mathbf{A}\hat{\mathbf{x}} = \mathbf{y} \Rightarrow \mathbf{A}(\hat{\mathbf{x}} - \mathbf{x}_o) = \mathbf{A}\hat{\mathbf{x}} - \mathbf{A}\mathbf{x}_o = \mathbf{y} - \mathbf{y} = \mathbf{0}.$

# Minimizing the $\ell^0$ Norm: Exhaustive Search

## Algorithm ( $\ell^0$ Minimization via Exhaustive Search):

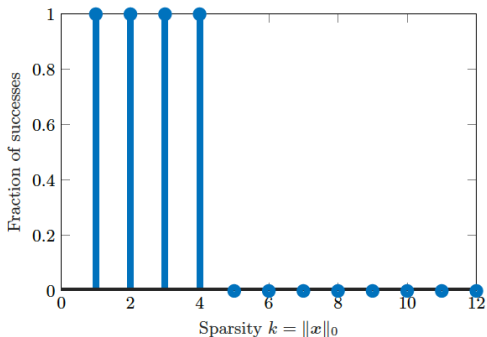
- 1: **Input:** a matrix  $\mathbf{A} \in \mathbb{R}^{m \times n}$  and a vector  $\mathbf{y} \in \mathbb{R}^m$ .
- 2: **for**  $k = 0, 1, 2, \dots, n$ ,
- 3:   **for** each  $I \subseteq \{1, \dots, n\}$  of size  $k$ ,
- 4:     **if** the system of equations  $\mathbf{A}_I \mathbf{z} = \mathbf{y}$  has a solution  $\mathbf{z}$ ,
- 5:       set  $\mathbf{x}_I = \mathbf{z}$ ,  $\mathbf{x}_{I^c} = \mathbf{0}$ .
- 6:     **return**  $\mathbf{x}$ .
- 7:   **end if**
- 8: **end for**
- 9: **end for**

**Example:** computational complexity in finding the sparsest solution via the exhaustive search algorithm.

# Minimizing the $\ell^0$ Norm: Simulations

$$\textbf{Solve: } \min \|x\|_0 \quad \text{s.t.} \quad \mathbf{A}x = y. \quad (28)$$

$\mathbf{A}$  is of size  $5 \times 12$ . Fraction of success across 100 trials.

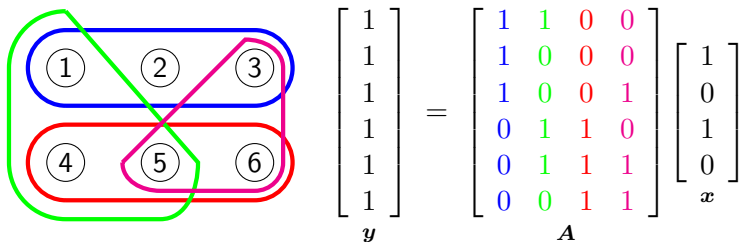


# Computational Complexity of $\ell^0$ Minimization

## Theorem (Hardness of $\ell^0$ Minimization)

The  $\ell^0$ -minimization problem  $\min \|x\|_0$  s.t.  $Ax = y$  is (strongly) **NP-hard**.

**Proof:** Reducible to an *Exact 3-Set Cover* (E3C) problem.



**Figure: Exact 3-Set Cover as a Sparse Representation Problem.** **Left:** a universe  $S = \{1, \dots, 6\}$  and four subsets  $U_1, \dots, U_4 \subseteq S$ .  $\{U_1, U_3\}$  is an exact 3-set cover. **Right:** the same problem as a linear system of equations. The columns of  $A$  are the incidence vectors for the sets  $U_1, U_2, U_3, U_4$ . The Exact 3-Cover  $\{U_1, U_3\}$  corresponds to a solution  $x$  to the system  $Ax = y$  with only  $m/3 = 2$  nonzero entries.



# Computational Complexity of $\ell^p$ Minimization ( $0 < p < 1$ )

## Theorem (Hardness of $\ell^p$ Minimization)

For any fixed  $0 < p < 1$ , the  $\ell^p$ -minimization problem

$$\min \|x\|_p^p \quad \text{subject to} \quad Ax = y$$

is (strongly) **NP-hard**.

## Corollary (Hardness of Smoothed $\ell^p$ Minimization)

For any fixed  $0 < p < 1$  and  $\epsilon > 0$ , the smoothed  $\ell^p$ -minimization problem

$$\min \|x + \epsilon\|_p^p \text{ s.t. } Ax = y \text{ is (strongly) } \mathbf{NP-hard}.$$

Nevertheless, computing a local minimizer of the problem can be done in polynomial time. (... **what about**  $\ell^1$ ?)

*A Note on the Complexity of  $L_p$  Minimization*, D. Ge, X. Jiang, and Y. Ye,  
*Mathematical Programming*, vol. 129, 2011.

## Two Fundamental Questions

To find the correct  $k$ -sparse solution to the under-determined linear system:

$$\mathbf{y} = \underset{m \times n, m \ll n}{\mathbf{A}} \times \underset{k \text{ sparse}}{\mathbf{x}}, \quad (29)$$

we want to know:

- ① **sample complexity**: how many measurements are needed for the problem to become computationally tractable? (**Part I**)
- ② **computational complexity**: once tractable, what is the precise computational complexity in finding the correct solution? (**Part II**)

*"It is quite probable that our mathematical insights and understandings are often used to achieve things that could in principle also be achieved computationally – but where blind computation without much insight may turn out to be so inefficient that it is unworkable."*

– Roger Penrose, *Shadows of the Mind*

# Assignments

- Reading: Section 2.1 and 2.2 of Chapter 2.
- Reading: Appendix B.
- Written Homework # 1.